

INTRODUCTION TO WQO THEORY

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Def: \mathcal{Q} is quasi-order \equiv reflexive & transitive

• $x \equiv y$ if $x \leq y$ & $y \leq x$

• $x < y$ if $x \leq y$ & $x \not\equiv y$

identifying \equiv -equiv elements produces a partial order

• x is minimal if $\forall y \leq x$ we have $y \equiv x$

Def: \mathcal{Q} is WQO $\equiv \forall$ infinite sequence $x_1, x_2, x_3, \dots \exists i < j$ s.t. $x_i \leq x_j$

Theorem: The following are equivalent

\hookrightarrow good / bad sequence

① \mathcal{Q} is WQO

② \forall sequence x_1, x_2, x_3, \dots contains an infinite non-decreasing subsequence $x_{i_1} \leq x_{i_2} \leq \dots$

③ \mathcal{Q} admits no infinite decreasing chains and no infinite antichains (non-equiv)

④ $\forall \emptyset \neq A \subseteq \mathcal{Q}$ contains at least 1, but only finitely many min. elements

Proof: ② \Rightarrow ①, ① \Rightarrow ③, ③ \Rightarrow ② by Ramsey's Theorem

③ \Rightarrow ④ no infinite decreasing chains $\Rightarrow \exists$ minimal element
no infinite antichains \Rightarrow only finitely many

④ \Rightarrow ③ by the same reasoning

Remark: If \aleph is a regular cardinal, and \mathcal{Q} WQO, then

$\forall \mathcal{Q}$ -sequence of length \aleph contains a non-decreasing subsequence of the same length

Def: A property $\mathcal{P}(x)$ is monotone $\equiv \forall x: \mathcal{P}(x) \& y \leq x \Rightarrow \mathcal{P}(y)$

• finitely testable $\equiv \exists$ finite set \mathcal{F} s.t. $\mathcal{P}(x) \Leftrightarrow (\forall F \in \mathcal{F}): F \not\leq x$

Theorem: \mathcal{Q} WQO $\Rightarrow \forall$ monotone property is finitely testable.

Proof: Define $\bar{\mathcal{Q}} = \{x \in \mathcal{Q} \mid \neg \mathcal{P}(x)\}$ and let \mathcal{F} be non-equiv min. elements of $\bar{\mathcal{Q}}$
 F_1, F_2, \dots, F_m

• if $\mathcal{P}(x)$ then $\forall F_i \not\leq x \dots$ if $F_j \leq x$ then $\mathcal{P}(F_j) \because \mathcal{P}$ is monotone $\&$

• if $\forall F_i \not\leq x$ then $\mathcal{P}(x) \dots$ if $\neg \mathcal{P}(x)$, take $\mathcal{Y} := \{y \in \bar{\mathcal{Q}} \mid y \leq x\}$

$\rightarrow \mathcal{Y}$ has a min. element $H \leq x$, and $H \equiv F_j$ for some $j \because H$ min. in $\bar{\mathcal{Q}}$

$\Rightarrow F_j \leq x \&$

Note: Graph minor theorem $\Rightarrow \forall$ surface \exists finite obstruction set

Sphere $\dots K_5, K_{3,3}$, Projective plane 35 obs., other than that not known

Preservation properties

- Q well-ordered $\Rightarrow Q$ WQO ... for example ω or $\{1, 2, \dots, n\}$
- Q finite $\Rightarrow (Q, =)$ WQO ... finite alphabets

Def: Product order of $P \times Q$: $(p, q) \leq (p', q')$ if $p \leq_P p'$ & $q \leq_Q q'$

Dickson's lemma: P, Q WQOs $\Rightarrow P \times Q$ WQO.

Proof: $(p_n, q_n)_n$, take ω -nondecreasing subsequence $p_{i_1} \leq p_{i_2} \leq \dots$
 $\Rightarrow (p_{i_n}, q_{i_n})_n$, Q WQO ... $\exists l < k$ s.t. $q_{i_l} \leq q_{i_k}$ ▣

Def: $Q^{<\omega}$ ordered by $(a_i)_{i < \omega} \leq (b_j)_{j < \omega}$ if \exists increasing $f: \mathbb{N} \rightarrow \mathbb{N}$ s.t.

- $(\Sigma, =)$ finite alphabet, $u \in v$ is word embedding $\forall i < n: a_i \leq b_{f(i)}$
- $P(Q)$ ordered by $X \leq_s Y$ if \exists inj. $f: X \rightarrow Y$ s.t. $\forall x \in X: x \leq f(x)$.

Theorem (Higman, 1952): Q WQO $\Rightarrow Q^{<\omega}$ WQO.

Proof: Construct a minimal bad sequence w_0, w_1, w_2, \dots min $|w_i|$

- write $w_i = a_i w'_i$
- Q WQO $\Rightarrow \exists$ increasing $f: \mathbb{N} \rightarrow \mathbb{N}$ s.t. $a_{f(0)} \leq a_{f(1)} \leq a_{f(2)} \leq \dots$

claim: $w'_{f(0)}, w'_{f(1)}, w'_{f(2)}, \dots$ is good

\hookrightarrow then $\exists i < j: w'_{f(i)} \leq w'_{f(j)}$ so $w_{f(i)} \leq w_{f(j)}$

Proof: Suppose not; then the sequence

$w_0, w_1, \dots, w_{f(0)-1}, w'_{f(0)}, w'_{f(1)}, \dots$ is bad

• if $w_j \leq w'_{f(k)}$ for some $j \leq f(0) \leq f(k)$ then $w_j \leq w_{f(k)}$ ∇

Hence it is bad, contradicting the minimality of $w_{f(0)}$ since $|w'_{f(0)}| < |w_{f(0)}|$ ▣

Kruskal's Tree Theorem 1960

Def: Order-Theoretic tree is a partial order (T, \leq) s.t. $\forall x \in T$,

$(\leftarrow, x) := \{y \in T \mid y < x\}$ is a finite chain

There is a unique minimal element ... root

Intuition: $y < x$ means "y lies on the unique path from x to the root"

Def: A homeomorphic embed of two \mathcal{Q} -labeled trees: $\varphi: S \rightarrow T$

- i) $\ell_S(x) \leq \ell_T(\varphi(x))$... respects labels Write $S \leq T$
 - ii) $x < y \Rightarrow \varphi(x) < \varphi(y)$
 - iii) $\varphi(x \vee y) = \varphi(x) \wedge \varphi(y)$
- } topological minor for rooted trees

Theorem (Kruskal): \mathcal{Q} WQO $\Rightarrow T_f(\mathcal{Q}) =$ finite \mathcal{Q} labeled trees WQO by \leq

Corollary: Finite graph-trees are WQO by topological minors ... Take $\mathcal{Q} = \{0, 1\}$
By Higman's lemma also forests.

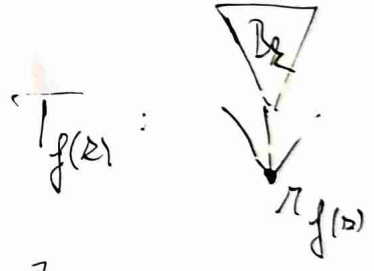
Proof: Construct a minimal bad sequence T_0, T_1, T_2, \dots minimizing $|T_i|$

- π_i ... root of T_i
- B_i ... branch-trees of T_i



claim: $B := \bigcup_{i < \omega} B_i$ is WQO by \leq

proof: Suppose that B_0, B_1, \dots is bad



- $\forall j$ select $f(j)$ or that $B_j \in B_{f(j)}$
- let k minimize $f(k)$
- consider the sequence $T_0, T_1, \dots, T_{f(k)-1}, B_k, B_{k+1}, \dots$

⊗ it is bad, contradicting the minimality of $T_{f(k)}$ since $|B_k| < |T_{f(k)}|$

↳ if $T_i \leq B_j$ for some $i < f(k)$ and $j \geq k$ then $T_i \leq T_{f(j)}$ & $i < f(k) \leq f(j)$

Because \mathcal{Q} is WQO, we may assume that $\ell(\pi_0) \leq \ell(\pi_1) \leq \dots$

Since B is WQO, Higman's lemma gives (B, \leq_s) WQO, so

$\exists i < j$ s.t. $B_i \leq_s B_j$

Together, $T_i \leq T_j$, contradicting badness of the minimal bad sequence

Ding's Theorem 1992

cycles \diamond, \square, \dots and forks $F_n \rightsquigarrow, \rightsquigarrow$ are induced subgraph \subseteq_i and subgraph \subseteq antichains

Theorem (Ding) If \mathcal{G} is a \subseteq -closed class of finite graphs, then the following are equiv:

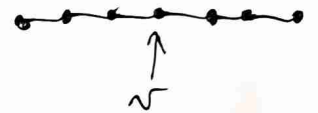
- ① \mathcal{G} is WQO by \subseteq
 - ② \mathcal{G} is WQO by \subseteq_i
 - ③ \mathcal{G} contains only finitely many C_n and F_n
- clearly:
② \Rightarrow ① \Rightarrow ③

Theorem (Ding) If $\exists m$ s.t. no $G \in \mathcal{G}$ has $P_m \subseteq G$, then \mathcal{G} is WQO by \subseteq_i .

Def: Tree-depth of a graph G is

$$td(G) := \begin{cases} 1, & \text{if } |G|=1 \\ \max_i td(G_i), & \text{if } G \text{ is disconnected with components } G_1, \dots, G_m \\ 1 + \min_{v \in G} td(G-v), & \text{otherwise} \end{cases}$$

① $td(P_m) \approx \log_2(m) \dots td(P_m) = 1 + td(P_{\lfloor m/2 \rfloor})$



\rightarrow if $td(G) < k$ then $P_{2^k} \not\subseteq G$ because td is minor monotone

② if $P_m \not\subseteq G$, then $td(G) < m$

\hookrightarrow if $td(G) \geq m$, then by taking the best v in \textcircled{X} we construct path $v_1, \dots, v_m \subseteq G$

Def: Induced subgraphs for \mathcal{Q} -labeled graphs

$$(G, l_G) \subseteq_i (H, l_H) \text{ if } \exists \text{ injection } f: V_G \rightarrow V_H \text{ s.t.}$$

- ① $uv \in E_G \Leftrightarrow f(u)f(v) \in E_H$
- ② $\forall v: l_G(v) \leq l_H(f(v))$

Notion: $D_m(\mathcal{Q}) \dots$ finite \mathcal{Q} -labeled graphs with tree-depth $\leq m$

Theorem: $\forall m: \mathcal{Q}$ WQO $\Rightarrow D_m(\mathcal{Q})$ WQO by \subseteq_i .

Proof: Induction on $m \dots$ case $m=1$ holds from Higman's lemma

- let $m > 1$, sequence $(G_i, l_i) \dots$ WLOG $\forall i: td(G_i) \geq 2 \dots$ there are ∞ many i
- $\mathcal{Q}^+ := \mathcal{Q} \times \{0,1\}$ ordered by $(q,e) \leq (q',e')$ if $q \leq q'$ & $e=e'$ is WQO
- $\forall i$ pick $v_i \in G_i$ s.t. $td(G_i - v_i) \leq m-1 \dots H_i := G_i - v_i$
- since \mathcal{Q} is WQO we may assume that $l_0(v_0) \leq l_1(v_1) \leq \dots$
- \mathcal{Q}^+ -label H_i by $l_i^+(u) := (l_i(u), e_i(u))$ where $e_i(u)$ encodes whether $uv_i \in E(G_i)$
- Induction hypothesis: $D_{m-1}(\mathcal{Q}^+)$ is WQO $\Rightarrow \exists i < j$ s.t. $(H_i, l_i^+) \subseteq_i (H_j, l_j^+)$
- let $\varphi: V(H_i) \rightarrow V(H_j)$ witness this, then $\varphi \cup \{(v_i, v_j)\}$ witnesses $(G_i, l_i) \subseteq_i (G_j, l_j)$

Excursion to infinity

Rado's counterexample: Higman's lemma fails for \mathbb{Q}^ω

1954

$\mathbb{Q} = \{(i, j) \mid i < j < \omega\}$ ordered by $(i_1, j_1) \leq (i_2, j_2)$ if $i_1 = i_2 \ \& \ j_1 \leq j_2$ or $j_1 < i_2$

→ easy to verify by Dilworth's lemma that \mathbb{Q} is WQO

	1	2	3	4	5	6	...
0	(0,1)	(0,2)	(0,3)	(0,4)	(0,5)	(0,6)	...
1	(1,1)	(1,2)	(1,3)	(1,4)	(1,5)	(1,6)	...
2		(2,3)	(2,4)	(2,5)	(2,6)	...	
3			(3,4)	(3,5)	(3,6)	...	
4				(4,5)	(4,6)	...	

$P(\mathbb{Q})$ is not WQO

by order-respecting maps

If ~~(1,2)~~ \in ~~(2,3)~~, then

$\circ \leq$ something green

Better-quasi-orderings ... Nash-Williams 1965

\mathbb{Q} BQO $\Rightarrow P(\mathbb{Q}), \mathbb{Q}^{<\omega}$ BQO ... $\mathbb{Q}^{<\omega} =$ transfinite sequences

\mathbb{Q} BQO $\Rightarrow T_\omega(\mathbb{Q}) =$ finite or infinite \mathbb{Q} -labeled trees BQO by \leq

↳ thus the class of all graph-trees is WQO by top-minors

\mathbb{Q} BQO $\Rightarrow D_n^{inf}(\mathbb{Q}) =$ graphs of tree-depth $\leq n$ BQO by \leq_i

Some statements do not need BQO theory

\mathbb{Q} WQO \Rightarrow transfinite sequences $\mathbb{Q}^{<\omega}$ with finite range WQO

Def: $S \leq_a T$ if \exists mapping $f: S \rightarrow T$ preserving immediate successors

↳ tree-homorphism

Theorem (Smolík, 2026): One can prove without BQO theory that

$T_\omega =$ finite or infinite trees are WQO by tree-homomorphisms.